#### Math 4803

### March 25, 2024

#### LAST TIME

Bounded polygons in (R², deuc), (H², dhyp), {(S², dsph) are <u>Compact</u>, and thus their quotients by edge gluings are <u>complete</u>.

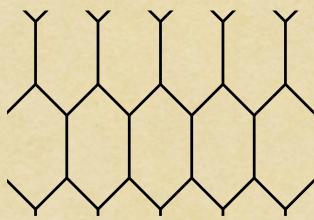
Upshot: Edge gluings of bounded polygons which satisfy the angle conditions give tessellations.

#### **TODAY**

- · Precluding some tessellations of (TR2, denc).
- · An example in (HT, day).
- · Tessellations by triangles.
- · A tessellation by an unbounded polygon.

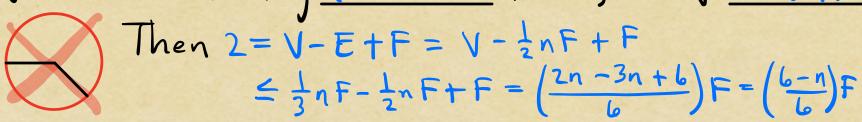
## Convex tilings of (R2, deuc)

tact. There is no tessellation of (R, deuc) whose tiles are convex polygons with more than lesides. (Fake proof/reason it's true.)



A tessellation of (R, deuc) gives us a Planar graph (kind of), and these Satisfy Euler's formula; V-E+F=Z

Each edge is shared by 2 faces, So E = \frac{1}{2}nF, where each face has nedges. Convexity ensures that each vertex is shared by at least 3 faces, so  $V \leq \frac{1}{3}nF$ .

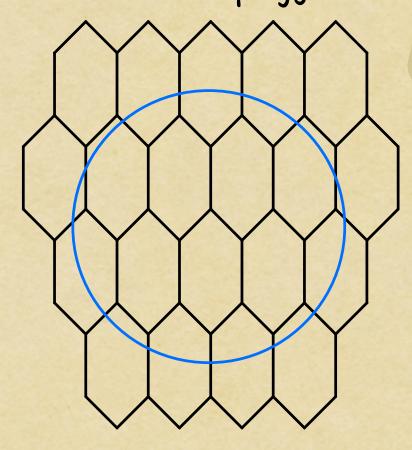


It n>7, then the RHS is negative. X



## Convex tilings of (R2, deuc)

Fact. There is no tessellation of (IR2, deuc) whose tiles are convex polygons with more than lesides.



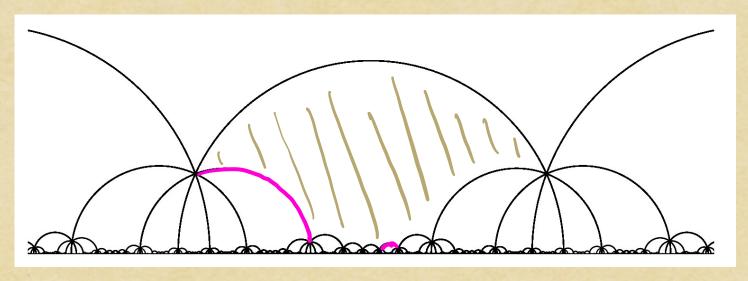
Problem: We don't really have a <u>Planar graph</u>.

Fix: For R>0, just think about the faces intersecting Bdenc (Po,R) Then F~ R<sup>2</sup>

Comparing E & V to Fgets more annoying, but it works out.

F<sub>2</sub>~R.

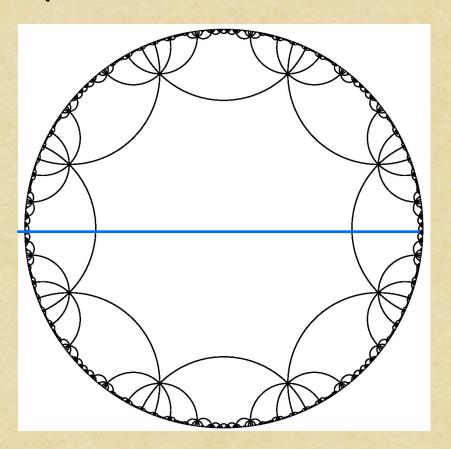
Tessallation (s) of (H², dhyp) by bounded octagons
Unsurprisingly, (H², dhyp) admits some tessellations which are prohibited in (R², deuc). For example, we built an octagon X CH² which admits an edge gluing by translations
We can use this to tessellate (H², dhyp):

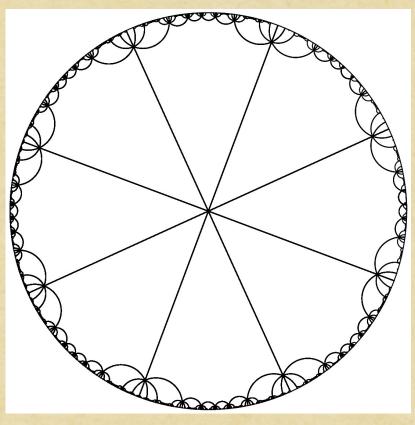


 $F_1 \sim 4\pi \sinh^2(R/2)$  $F_3 \sim 2\pi \sinh(R)$ 

Why doesn't our Euler's formula argument work? F\_ ~ F\_ , so we can't ignore effect of frontier tiles

# Tessallation (s) of (Hi, dhyp) by bounded octagons This one looks even cooler in the disc model:





By tessellating (H², dny) with lots of different triangles (coming up), we can construct tessellations by a variety of polygons.

Tessellations by triangles
Thm. For any integers a, b, c>2, (1) if  $\frac{\pi}{a} + \frac{\pi}{b} + \frac{\pi}{c} = \pi$ , there is a tessellation of (R, denc) by Euclidean triangles of angles =, =; (2) if I+ I+ I < TT, there is a tessellation of (H, days) by hyperbolic triangles of angles =, =; (3) if I+ I+ I>TT, there is a tessellation of (S, dsph) by Spherical triangles of angles =, =, =. (Proof.) Prop 5.13 => the triangle we want exists. For gluing data we take Then the vertices fall into 3 equivalence classes, and their angle sums are Tisa. Since T is a bounded polygon, our previous theorems give a tessellation. RMK. Our previous proof in (R, denc) used a different gluing.

Tessellations by triangles How many tessellations does this give us?

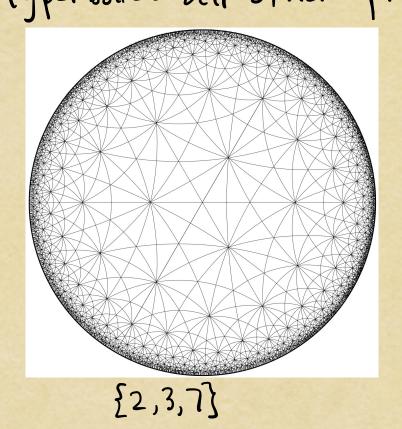
Fudidean: \$ ±+ ±=π, with a≤b≤c.

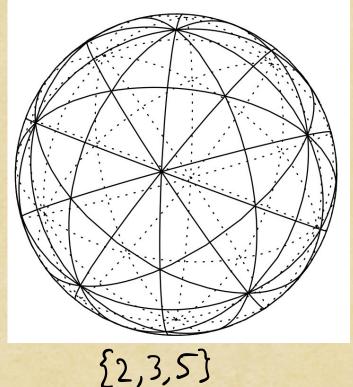
Then  $\frac{\pi}{a} \Rightarrow \frac{\pi}{3} \rightarrow a = 2$  or a = 3If a = 2, then  $\frac{\pi}{b} + \frac{\pi}{c} = \frac{\pi}{2} \rightarrow \frac{\pi}{b} \Rightarrow \frac{\pi}{4} \rightarrow b = 3$  or b = 4If b = 3, then  $\frac{\pi}{c} = \frac{\pi}{2} - \frac{\pi}{3} = \frac{\pi}{6} \rightarrow \{2,3,6\}$ If b = 4, then  $\frac{\pi}{c} = \frac{\pi}{2} - \frac{\pi}{4} = \frac{\pi}{4} \rightarrow \{2,4,4\}$ If a = 3, then  $\frac{\pi}{b} + \frac{\pi}{c} = \frac{2\pi}{3} \Rightarrow b = c = 3$ 

{3,3,3}

Tessellations by triangles How many tessellations does this give us? Endidean: {a,b,c}={2,3,6},{2,4,4}, or {3,3,3}. Spherical: \$ II + II + II > II, with a \le b \le C. Then a < 3. If a=2, then T+ = > T = > 6 < 4 If b=2, then  $\frac{\pi}{c} > 0 \rightarrow \{2,2,c\}$ If b=3, then  $\frac{\pi}{c} > \frac{\pi}{b} \rightarrow c < b$ This gives {2,3,3}, {2,3,4}, {2,3,5}

Tessellations by triangles How many tessellations does this give us? Endidean: [a,b,c]=[2,3,6],[2,4,4], or [3,3,3]. Spherical: {a,b,c}={2,3,3},{2,3,4},{2,3,5}, or {2,2,c}. Hyperbolic: all other options



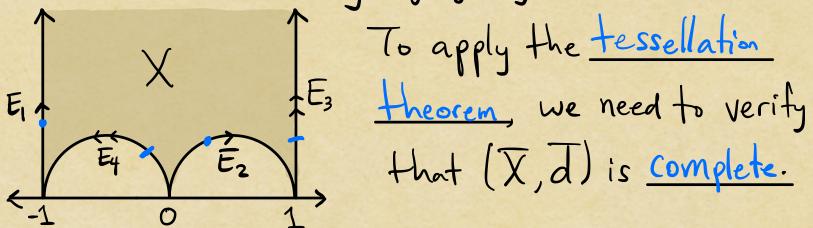


# At essellation of (H2, dhyp) by an unbounded square All of today's tessellations use the following:

Prop. Let X be a bounded polygon in (IR, deuc), (H, dhyp), or (S, dsph). Then any edge gluing of X which satisfies the angle condition on vertices leads to a tessellation.

This works because edge gluings on bounded polygons lead to <u>complete</u> quotient spaces. Now we'll see an example which uses an unbounded polygon.

Consider the following edge gluing of a square:



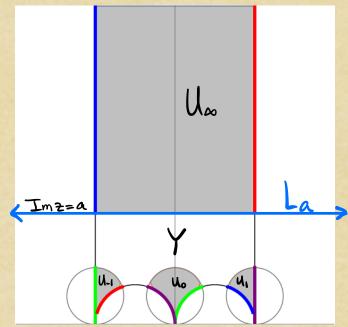
## A tessellation of (Hi, dhyp) by an unbounded square

Recall that we can decompose X as

$$X = U \circ Y$$

with Y a bounded polygon and

where, for instance,



This gives a decomposition

$$\overline{X} = \overline{V} \circ \overline{Y}$$
,

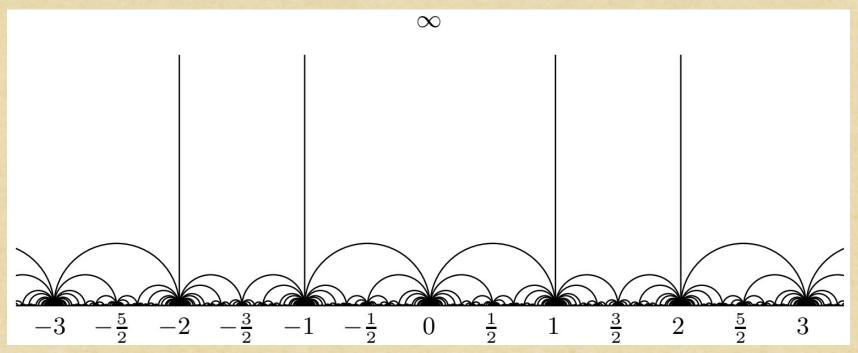
and X is complete iff

U and Y are complete. Note that Y a bounded polygon"

⇒ T is complete. It remains to check U.

A tessellation of (H2, dhyp) by an unbounded square Claim. U is complete. First, we know that U is isometric to the Pseudosphere, (Sa, dsa) C (R, denc). Be cause (TR, deuc) is complete, any finite-length sequence in (Sa, dsa) will converge in (R, denc). But Sa is determined by the equations and inequalities X = t - tanht,  $y = secht \cdot coss$ ,  $z = secht \cdot sins$ , No stiet cosh'  $(a\pi/3) \le t \le \infty$ ,  $0 \le s \le 2\pi$ . inequalities and thus the limit point will be contained in Sa. So any finite-length sequence in (Sa, dsa) will converge in (Sadsa). Sa is complete > U is complete > X is complete -> we get a tessellation

## A tessellation of (Hidny) by an unbounded square



The edge gluing is not by translation!