# Math 4803

#### **February 7, 2024**

RECENTLY

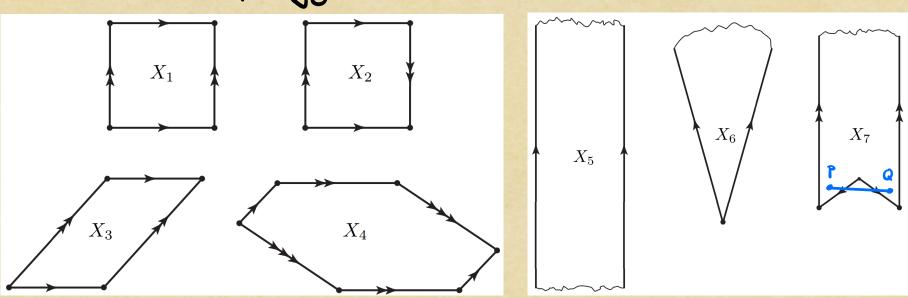
Quotients of metric spaces, which are generally  $\frac{\text{Semi-metric}}{\text{J}(P,Q)=0} \Rightarrow P=Q$ 

#### **TODAY**

Quotients of polygons via edge gluings.

Euclidean polygons A polygon in (R, deuc) is a region XCR whose bdry de composes as  $\partial X = E_1 \cup E_2 \cup \cdots \cup E_n$ , where (1) each Ei is a line segment, line, or half-line; (2) each nonempty Ein Ei is an endpoint of Eig Ei, for itj; (3) for distinct  $1 \le i,j,k \le n$ ,  $EinEjnEk = \emptyset$ . We call each Ei an edge of X and Call the nonempty Ein Ej vertices. We require X to be closed, meaning that dx CX, but do not require X to be bounded.

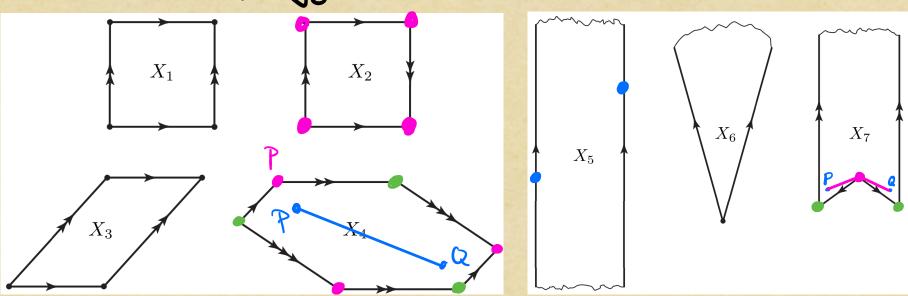
# Euclidean polygons



We call X bounded if 3 PER and r>0 s.t.

We call X convex if, YP,QEX, there is a geodesic arc contained in X connecting P to Q.

# Euclidean polygons



The Euclidean path metric on X is defined by  $d_x(P,Q) := \inf\{l_{denc}(Y) \mid P \stackrel{\sim}{\sim} Q \text{ in } X\}$ , for all  $P,Q \in X$ .

Check: If X is Convex, then  $d_X = d_{enel}_{X \times X}.$ 

Edge gluings An edge gluing of a Euclidean polygon is a collection of isometries (w.r.t. deue)  $f_{2k-1}: E_{2k-1} \longrightarrow E_{2k}$   $f_{2k}:= f_{2k-1}$  by which pairs  $\{E_1, E_2\}, \{E_3, E_4\}, \dots, \{E_{2k-1}, E_{2k}\}$  of its edges. Namely, the edges Ezk-1 ; Ezk must have the Same length. If these edges are finite, we Can encode 42k-1 by decorating Ezk-1 & Ezk with arrows. If the edges are lines, we'll turther need dots to determine 2k-1. (See previous figures.)

Edge gluings An edge gluing of X gives us a partition: (4304)(P) E4 (4,04,042)(P) · PEX in the interior  $\Rightarrow \overline{P} = \{P\}$ .  $E_1$   $\int_{E_2}$ • PEX in exactly one edge  $E_i$   $\Rightarrow \overline{P} = \{P, \Psi_i(P)\}$ • PEX a vertex  $\Rightarrow \overline{P} = \{(\Psi_i \circ \Psi_i) \cap (P) \mid (\Psi_i \circ \dots \circ \Psi_{i_1}) \cap (P) \in E_{i_3}, \Psi_j\}$ • PEX in exactly one edge  $E_i$   $\Rightarrow \overline{P} = \{(\Psi_i \circ \Psi_{i_{k-1}} \circ \dots \circ \Psi_{i_1}) \cap (P) \mid (\Psi_{i_{j-1}} \circ \dots \circ \Psi_{i_1}) \cap (P) \in E_{i_j}, \Psi_j\}$ 

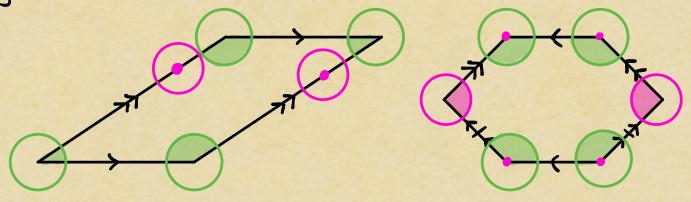
Thm. If X is obtained from the Euclidean polygon X by gluing together edge pairs by isometries, then the gluing is proper. (i.e., the semi-metric dx is a metric)

Euclidean surfaces
Our first example of a Euclidean surface was
Our first example of a Euclidean surface was the cylinder:  a  The cylinder:  The
We claimed that while distances between points will
Change thru our gluing process, the local geometre
will not. It's time to make this more precise.
Def. The metric space (X,d) is said to be locally isometric to (X',d') if, for every PEX, there exist
aball $B_{dx}(P, E)$ , a ball $B_{dx'}(Q, E)$ , and an

isometry between these.

#### Euclidean surfaces

We call (X,d) a <u>Enclidean Surface</u> if it is locally isometric to the Euclidean plane (IR², denc), and Say that d is a <u>Euclidean metric</u>.



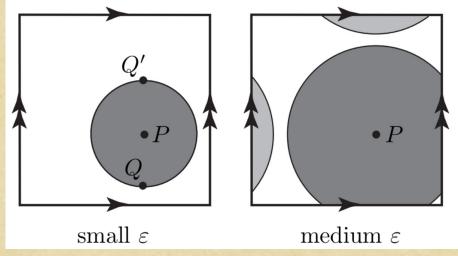
Thm. Let  $(X, d_X)$  be a Euclidean polygon, and let  $(\overline{X}, d_{\overline{X}})$  be the quotient <u>metric space</u> obtained from some edge gluing of  $(X, d_X)$ . Suppose that, for every vertex  $P \in X$ ,  $\sum_{Q \in \overline{P}} \Delta(Q) = 2\pi$ 

Then  $(X, d_{\overline{X}})$  is Euclidean surface

#### Proof interlude

Our two statements about Euclidean edge gluings — <u>Properness</u> and the criterion for <u>surfaces</u> — rely on two technical facts which we'll state without proof.

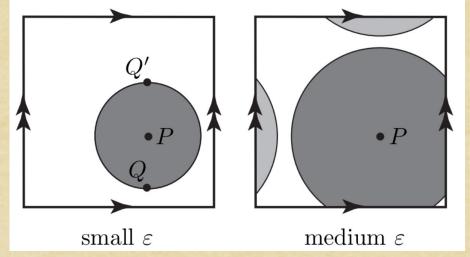
Lemma ① Let  $(X, \overline{d}_X)$  be obtained from the Euclidean Polygon (X, d) by an edge gluing. For every  $P \in X$ ,  $\exists \ \varepsilon_0 > 0$  s.t. for every  $\varepsilon \leq \varepsilon_0$  and every  $Q \in X$ ,



$$\overline{Q} \in \overline{B}_{dx}(\overline{P}, \varepsilon)$$
iff
 $\exists P' \in \overline{P} \text{ s.t.}$ 
 $Q \in \overline{B}_{dx}(P', \varepsilon)$ .

#### Proof interlude

Lemma () Let  $(X, d_X)$  be obtained from the Euclidean polygon (X, d) by an edge gluing. For every  $P \in X$ ,  $\exists \ E_0 > 0$  s.t. for every  $E \le E_0$  and every  $Q \in X$ ,



$$\overline{Q} \in \overline{B}_{dx}(\overline{P}, \varepsilon)$$
iff

 $\exists \overline{P}' \in \overline{P} \text{ s.t.}$ 
 $Q \in \overline{B}_{dx}(\overline{P}', \varepsilon).$ 

i.e., for suff. small  $\varepsilon > 0$ ,  $Ba_{\overline{X}}(\overline{P}, \varepsilon)$  looks like a union of balls in X

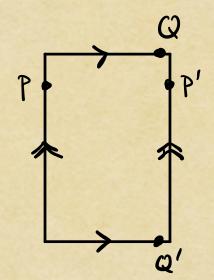
The proof involves a lot of <u>casework</u> using the <u>triangle</u> inequality.

### Proof of properness

We can now verify that  $d_X$  is a metric when  $(X, \overline{d})$  is obtained from a Euclidean polygon  $(X, \overline{d})$  via edge gluing.

 $P \neq Q$ . We NTS  $d_{x}(P,Q) > 0$ . Pick  $\varepsilon_{o} > 0$  for P as in Lemma (1).

 $P + \overline{Q} \Rightarrow \underline{P} \cap \overline{Q} = \emptyset$  as subsets of X



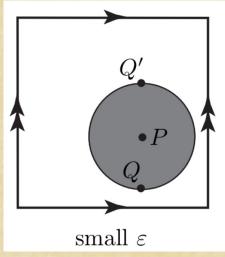
Q We can choose E<Eo s.t.

 $d_{x}(P,Q) > \varepsilon$ ,  $\forall P \in P \notin Q \in Q$ . (Relies on  $P \notin Q$  being finite.)

Then  $\overline{d}_{x}(\overline{P}, \overline{Q}) \geq \varepsilon$ .

#### Proof interlude

Our other theorem says that edge gluing a Euclidean polygon yields a Euclidean surface, provided the angle sums at the Vertices make



sense.

The proof is similar to that of Lemma 1, in that each PEX needs a ball small small ε enough for some purpose, and there's a lot of Casework.

The linchpin is

Lemma (2) Let 4:9 → g' be an isometry botwn two geodesics in (R2, d), and choose a "side" of each of 9,9'. There is a unique isometry  $\Psi:\mathbb{R}^2 \to \mathbb{R}^2$  which extends 4 and respects the sides.

# Hyperbolic surfaces

A hyperbolic polygon X < (Ht, day) is a region whose boundary in H consists of finitely many geodesics Meeting only at their endpoints. Exactly two edges

meet at each vertex. We say that X is <u>unbounded</u> if its

boundary includes a nontrivial interval on the real axis. \* or it's unbound unbounded

Gold region has trivial body on real axis > bounded.

txactly as in the Euclidean case, we can define the path metric dx on X, as well as edge gluings, Which provide pairwise isometries of the edges in days.

## Hyperbolic surfaces

Our two important theorems for edge gluings of Euclidean polygons also hold in the hyperbolic case:

Let (X, dx) be obtained from a hyperbolic polygon Via an edge gluing by hyperbolic isometries.

Thm. The semi-metric dx is in fact a metric.
i.e., the edge gluing is <u>proper</u>

Thm. Suppose that, for every vertex  $P \in X$ ,  $Z \times (Q) = 2\pi$ .

Then  $(X, d_{\overline{X}})$  is a hyperbolic surface. i.e.,  $(X, d_{\overline{X}})$  is locally isometric to  $(H^2, d_{Ny})$ 

#### Spherical surfaces

A spherical polygon  $X \subset (S^2, d_{sph})$  is a region whose boundary consists of finitely many geodesics Meeting only at their endpoints. Exactly two edges meet at each vertex. All spherical polygons are bounded.

We define the path metric and edge gluings as usual.

Finally, let  $(\overline{X}, \overline{d}_x)$  be obtained from a spherical polygon via an edge gluing by spherical isometries.

Thm. The semi-metric dx is in fact a metric.

Thm. Suppose that, for every vertex PEX,  $\sum_{Q \in P} L(Q) = 2\pi$ .

Then  $(X, d_X)$  is a Spherical surface. (s², dsph)

Coming up Lots of examples!